

# Impoverishment

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# Additional Operations in Halle & Marantz (1993)

- ▶ Impoverishment
- ▶ Fission
- ▶ Fusion
- ▶ Readjustment

# Impoverishment

Delete features of a syntactic head

$$\begin{bmatrix} +1 \\ +pl \end{bmatrix} \rightarrow [+1]$$

$$\begin{bmatrix} +1 \\ +pl \end{bmatrix} \rightarrow [ \ ]$$

# The Double Role of Impoverishment

- ▶ Impoverishment serves as a **repair** operation for cases where word( form)s get the “wrong” exponent/vocabulary item
- ▶ Impoverishment captures **generalizations** on syncretism which are independent of single vocabulary items

# Impoverishment as Repair

# Impoverishment as Repair

	sg	pl
<b>1</b>	leg- <b>e</b>	leg- <b>en</b>
<b>2</b>	leg- <b>st</b>	leg- <b>t</b>
<b>3</b>	leg- <b>t</b>	leg- <b>en</b>

	sg	pl
<b>1</b>	leg-t- <b>e</b>	leg-t- <b>en</b>
<b>2</b>	leg-t- <b>est</b>	leg-t- <b>et</b>
<b>3</b>	leg-t- <b>e</b>	leg-t- <b>en</b>

## Vocabulary Items

[+2 -pl]	↔	<b>st</b>
[-2 -1]	↔	<b>t</b>
[-2]	↔	<b>e</b>

## Problem:

Violation of the  
Subset Principle

# Impoverishment as Repair

**Syntax:** [+Agr -2 -1 -pl]

**Impoverishment:** [1] → ∅ / \_\_\_\_ [+past]

[+Agr -2 -pl]

## Vocabulary Items

[+2 -pl] ↔ **st**

[-2 -1] ↔ **t**

[-2] ↔ **e**



[-2] ↔ **e**

# Impoverishment as Generalization



# A System-wide Syncretism Pattern

	sg	pl
<b>Present</b>	1 leg- <b>e</b>	leg- <b>en</b>
	2 leg- <b>st</b>	leg- <b>t</b>
	3 leg- <b>t</b>	leg- <b>en</b>

	sg	pl
1	bi- <b>n</b>	sind- <b>Ø</b>
2	bi- <b>st</b>	sei- <b>t</b>
3	is- <b>t</b>	sind- <b>Ø</b>

	sg	pl
<b>Past</b>	1 leg-t- <b>e</b>	leg-t- <b>en</b>
	2 leg-t- <b>est</b>	leg-t- <b>et</b>
	3 leg-t- <b>e</b>	leg-t- <b>en</b>

	sg	pl
1	war- <b>Ø</b>	war- <b>en</b>
2	war- <b>st</b>	war- <b>t</b>
3	war- <b>Ø</b>	war- <b>en</b>

In the past tense 1sg forms are always identical to 3sg forms

# Capturing System-wide Syncretism by Impoverishment

$$+/-1 \rightarrow \emptyset / \text{---} [+Past]$$

	sg	pl
<b>1</b>	[+1 -2 -Pl]	[-1 -2 +Pl]
<b>2</b>	[-1 +2 -Pl]	[-1 +2 +Pl]
<b>3</b>	[-1 -2 -Pl]	[-1 -2 +Pl]

 $\Rightarrow$ 

	sg	pl
<b>1</b>	[ -2 -Pl]	[+1 -2 +Pl]
<b>2</b>	[-1 +2 -Pl]	[-1 +2 +Pl]
<b>3</b>	[ -2 -Pl]	[-1 -2 +Pl]

$\Rightarrow$  No vocabulary insertion can break the identity of 1sg and 3sg


# Impoverishment and Restrictiveness

# An Alternative to Impoverishment: Rules of Referral


Rules of Referral stipulate the identity of specific paradigm cells:

In the past tense 3sg verb forms are identical to corresponding 1sg verb forms

	<b>Singular</b>	<b>Plural</b>
<b>1</b>	legte	legten
<b>2</b>	legtest	legtet
<b>3</b>	legte	legten



	<b>Singular</b>	<b>Plural</b>
<b>1</b>	war	waren
<b>2</b>	warst	wart
<b>3</b>	war	waren



## Sind Rules of Referral Inhärent Paradigmatisch?

- ▶ Eine Art Rules of Referral zu verstehen ist als asymmetrische Verweise zwischen Paradigmenzellen
- ▶ In einer postsyntaktischen DM-artigen Architektur kann man sie aber auch als Regeln verstehen, die Merkmale vor Vocabulary Insertion verändern:

[+3-pl] → [+1-pl] / \_\_\_\_\_ [+past]

- ▶ Dann funktionieren Rules of Referral ähnlich wie (aber weniger restriktiv als) Impoverishment-Regeln

# A Simple Example Paradigm

$[-\text{masc } -\text{pl}]_1$	$[-\text{masc } +\text{pl}]_2$
$[\text{+masc } -\text{pl}]_3$	$[\text{+masc } +\text{pl}]_4$

 $\approx$ 

$a_1$	$b_2$
$c_3$	$d_4$

# Possible Types of Syncretism

**Type 0**

a	b
c	d

**Type 1**

a	b
---	---

**Type 2**

a	
	b

**Type 3**

a	b
b	a

# Deriving Type-0 Syncretism

	$[-m -p]$	$[-m +p]$	$[-+m -p]$	$[+m +p]$
<b>Insertion:</b>	a: $[-m -p]$	b: $[-m +p]$	c: $[-+m -p]$	d: $[+m +p]$

a	b
c	d



# Deriving Type-1 Syncretism

	$[-m -p]$	$[-m +p]$	$[+m -p]$	$[+m +p]$
<b>Impoverishment:</b>	$[-m \quad ]$	$[-m \quad ]$	$[+m \quad ]$	$[+m \quad ]$
<b>Insertion:</b>	$[-m]:a$		$[+m]:b$	

$$p \rightarrow \emptyset$$

a	b
---	---

# Deriving Type-2 Syncretism

	$[-m -p]$	$[-m +p]$	$[+m -p]$	$[+m +p]$
	$[-m -p]$	$[ \quad +p]$	$[+m -p]$	$[+m +p]$
	$[-m \quad ]$	$[ \quad \quad ]$	$[+m \quad ]$	$[+m \quad ]$
	$[-m \quad ]$	$[ \quad \quad ]$	$[ \quad \quad ]$	$[ \quad \quad ]$
	$[-m ]:a$	$[ \quad ]:b$		

$$-m \rightarrow \emptyset / [ \text{---} +p ]$$

$$p \rightarrow \emptyset$$

$$+m \rightarrow \emptyset$$

a	
	b

# Why Type-3 Syncretism Cannot Be Derived

$[-\text{masc } -\text{pl}]_1$	$[-\text{masc } +\text{pl}]_2$
$[+\text{masc } -\text{pl}]_3$	$[+\text{masc } +\text{pl}]_4$

 $\star \Rightarrow$ 

a	b
b	a

To derive this paradigm, the single cells must be impoverished such that:

- (i)  $\text{Cell}_1 = \text{Cell}_4$
- (ii)  $\text{Cell}_2 = \text{Cell}_3$
- (iii)  $\text{Cell}_{1,4} \neq \text{Cell}_{2,3}$

## Why Type-3 Syncretism cannot be Derived

$[-\text{masc } -\text{pl}]_1$	$[-\text{masc } +\text{pl}]_2$
$[+\text{masc } -\text{pl}]_3$	$[+\text{masc } +\text{pl}]_4$

 $\ast \Rightarrow$ 

<b>a</b>	<b>b</b>
<b>b</b>	<b>a</b>

The only way to guarantee that  $\text{Cell}_1 = \text{Cell}_4$   
is to impoverish both cells to [ ]

The only way to guarantee that  $\text{Cell}_2 = \text{Cell}_3$   
is to impoverish both cells to [ ]

but this results in complete syncretism for all 4 cells ( $\text{Cell}_{1,4} = \text{Cell}_{2,3}$ )