Infinity

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Infinite Candidate Sets: ba

	Delete none	**Delete a**	**Delete b**	**Delete both**
Insert None	ba	b	a	
Insert One	□ba b□a ba□	□b b□	□a a□	□
Insert Two	□□ba □b□a □ba□ b□□a b□a□ ba□□	□□b □b□ □b□ b□□ b□□ b□□	□□a □□a □a□ □□a □a□ a□□	□□
Insert Three	. . .	. . .	. . .	. . .

Computing Infinite Candidate Sets I

Initial Candidate Set = {a,b}*

0	∅
1	a, b
2	aa, ab, ba, bb
3	aaa, aab, aba, abb, baa, bab, bba, bbb
4	aaaa, aaab, aaba, aabb, abaa, abab, abba, abbb, baaa, . . .
	. . .

Computing Infinite Candidate Sets II

ONSET: Each a is preceded by at least one b

0	∅
1	b
2	ba, bb
3	bab, bba
4	babb, bbab, bbba, baba
	. . .

Computing Infinite Candidate Sets III

MAX: For each a,b in the input
there should be at least one identical sound in the output

0	
1	
2	ba
3	bab, bba
4	babb, bbab, bbba, baba
	. . .

Computing Infinite Candidate Sets IV

NOCODA: Each b immediately precedes an a

0	
1	
2	ba
3	
4	baba
	. . .

Computing Infinite Candidate Sets V

DEP: The output shouldn't contain more sounds than the input

0	
1	
2	ba
3	
4	